

Automatic Detection of Unsafe Dynamic Component Loadings in Multi-Terminals by Using IP Address

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Abstract: The mechanism widely used for an improved system of modularity and flexibility is Dynamic loading of software components. Dynamic loading can be hijacked by placing an arbitrary file with the specified name in a directory searched before resolving the target component, thereby making the system more vulnerable. The vulnerable components are important to be deducted.

In this paper an automated technique to detect vulnerable and unsafe dynamic component loadings are presented. Analysis has two phases:

- 1) Online phase – by applying dynamic binary instrumentation to collect runtime information on component loading.
- 2) Offline phase – to analyze the collected information to detect vulnerable component loadings.

This technique is implemented in networked system of Microsoft Windows and UNIX using system IP address and it can deduct unsafe components and stopped. Our evaluation results show that unsafe component loading is prevalent in software on both OS platforms, and it is more severe on Microsoft Windows.

Keywords: Full path, Filename, Online Phase, Offline phase

I. Introduction

Dynamic loading allows an application the flexibility to dynamically link a component and use its exported functionalities. Benefits include modularity and generic interfaces for third-party software such as plug-ins. It also helps to isolate software bugs as bug fixes of a shared library can be incorporated easily. With these advantages, dynamic loading is widely used in designing and implementing software.

Important stage in dynamic loading is component resolution by locating the correct component for use at runtime. Two resolution methods are used in Operating systems by either specifying the *full path* or the filename of the target component. Using *full path*, operating systems simply locate the target from the given full path. With filename, operating systems resolve the target by searching a sequence of directories, determined by the runtime directory search order, to find the first occurrence of the component.

Operating systems might provide mechanisms to protect system resources. For example, Microsoft Windows supports Windows Resource Protection (WRP) to avoid system files from being replaced. However, these do not prevent loading of a malicious component located in a directory searched before the directory where the intended component resides.

In case an attacker sends a victim an archive file containing a document for a vulnerable program (e.g., a Word document) and a malicious DLL. If the victim opens the document after extracting the archive file, the vulnerable program will load the malicious DLL, which leads to remote code execution.

II. Unsafe Component Loading

This section describes dynamic loading of components, types of unsafe loadings and remote attack for vulnerable dynamic loading.

2.1 Dynamic Loading of Components

Software components often use functionalities exported by other components such as shared libraries at runtime. This operation is generally involves in three phases: **resolution, loading, and usage**. Specifically, an application resolves the needed target components, loads them, and utilizes the desired functions provided by them.

Component interoperation can be achieved through dynamic loading provided by operating systems or runtime environments. For instance, the **LoadLibrary** and **dlopen** system calls are used for dynamic loading on Microsoft Windows and Unix-like operating systems, respectively. Dynamic loading is mostly executed in two steps: **component resolution** and **chained component loading**.

2.1.1 Component Resolution

To resolve a target component, it is required to specify it correctly. Operating systems provide two types of target component specifications: **full path** and **filename**. For full path specification, operating systems resolve a target component based on the provided full path. For filename specification, operating systems obtain the full path of the target component from the provided file name and a dynamically determined sequence of search directories. Specifically, an operating system iterates through the directories until it finds a file with the specified file name, which is the resolved component. For example, suppose that a target component is specified as `audmigplugin.dll` and the directory search order is given as

C:\Program Files\iTunes; C:\Windows\System32; ...; \$PATH on Microsoft Windows. If the first directory containing a file with the name audmigplugin.dll is C:\Windows\System32, the resolved full path is determined by this directory.

2.1.2 Chained Component Loading:

In dynamic loading, the full path of the target component is determined by its specification through the resolution process and the component is incorporated into the host software if it is not already loaded. During the process of incorporation the target component, the component’s load-time dependent components are also loaded.

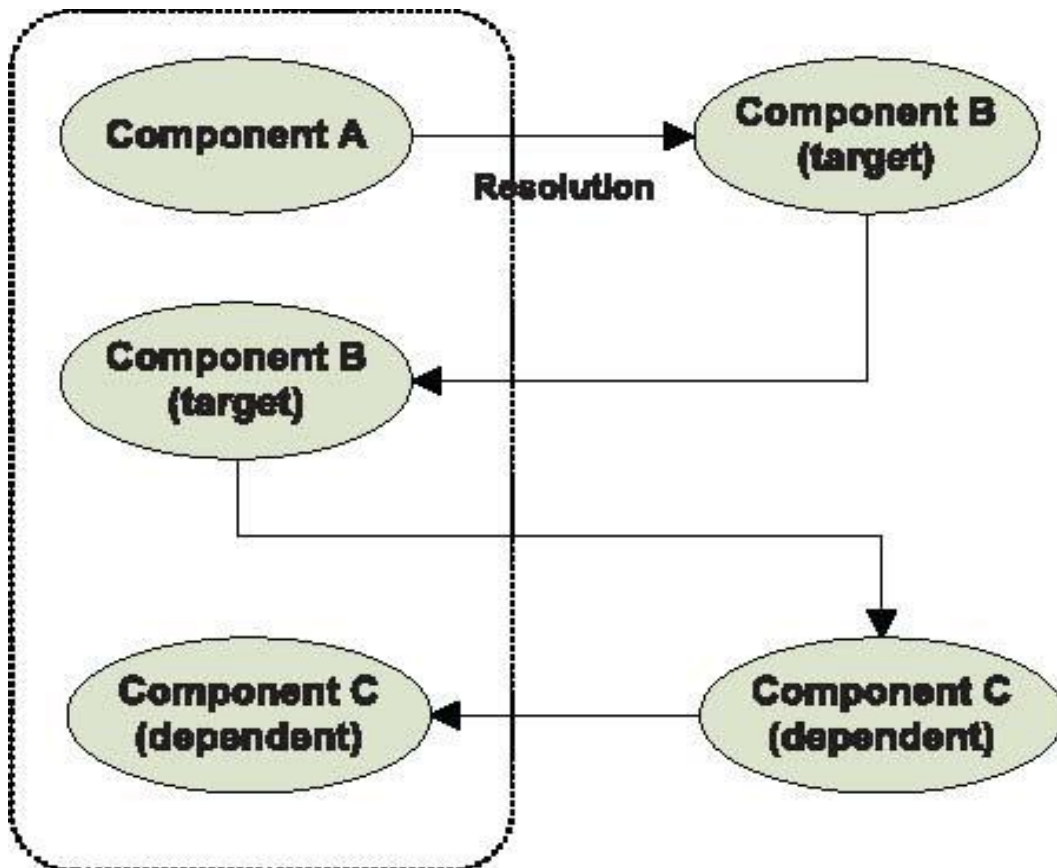


Fig. a. Dynamic component loading procedure

Fig. a, demonstrates the general procedure of dynamic loading. Suppose component B is loaded by component A. B’s dependent components (component C) are also loaded. We can usually obtain information on B’s dependent components from B’s file description. This process of chained component loading is repeated until all dependent components are loaded.

2.2 Unsafe Component Resolution

Despite dynamic loading is a critical step in software execution, it also has an inherent security implication. Particularly, a loaded target component is only determined by the specified file name. This can lead to the loading of unintended or even malicious components and thus may allow arbitrary code execution. Two types of unsafe component resolution are classified, **resolution failure** and **unsafe resolution**, and illustrated their conditions in Table 1.

TABLE 1: Conditions for Detecting Unsafe Component Loadings

Type	Conditions
Resolution failure	<ul style="list-style-type: none"> • Target component not found
Unsafe resolution	<ul style="list-style-type: none"> • the target component specified by its file name, • the resolution is determined by iteratively searching a sequence of directories, and • There exists another directory searched before the one containing the target component.

III. Detection of Unsafe Components in Network System Fig 2 - ARCHITECTURE DIAGRAM

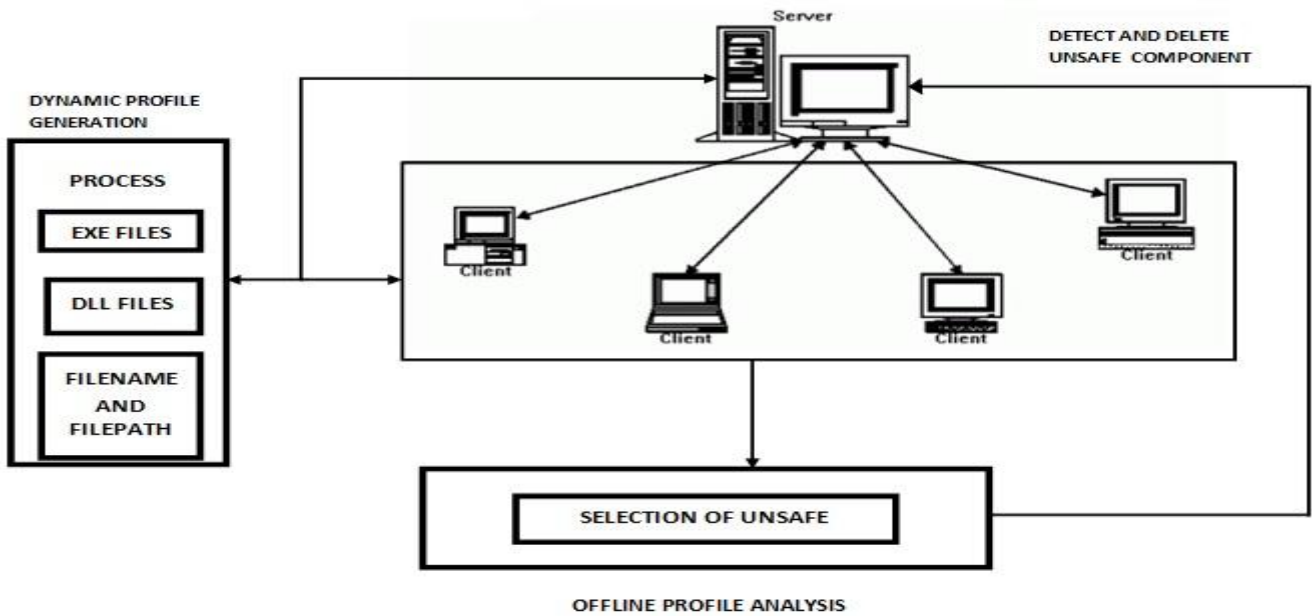


Fig. 2 shows the high-level overview of our analysis process, which is composed of two phases: dynamic profile generation and offline profile analysis.

3.1 Dynamic Profile Generation

We dynamically applied the binary executable analysis to capture a sequence of system level actions for dynamic loading of components. During the instrumented program execution, we collect three types of information: **system calls invoked for dynamic loading, image loading, and process and thread identifiers**. The collected information is stored as a profile for the instrumented application and is analyzed in the offline profile analysis phase.

The dynamic profile generated information extracted from the Client side to monitor the server by using the IP address. The server can monitor the dynamic loading to detection unsafe components and stop the vulnerable host.

3.2 Offline Profile Analysis

Firstly, we extract each component loading from the profile. First group a sequence of actions in the profile by process and thread identifiers as the actions performed by different threads may be interleaved due to context switching. This grouping separates the sequences of dynamic loadings performed by different threads. Then, we divide the sequence for each thread into subsequence of actions, one for each distinct dynamic loading. This can be achieved by using the first invoked system call for dynamic loading as a delimiter. Next, we obtain a list of groups, each of which contains a sequence of actions for loading a component at runtime. This gives the possible control flows in the dynamic loading procedure. Note that each group contains loading actions for both the target component and the load-time dependent components

Algorithm 1. OfflineProfileAnalysis

Input: S (a sequence of actions for a dynamic loading) Auxiliary functions:

Target Spec(S): return target specification of S

DirSearchOrder(S): return directory search order used

In S

ImgLoad(S): return the image loadings in S

Resolution Failure(S): return the resolution failures in S

Chained Loading(S): return actions for the chained

Loadings in S IsUnsafeResolution (filename, resolved path, search_dirs): check whether the resolution is unsafe

1: img_loads ImgLoad(S)

2: failed_resolutions Resolution Failure(S)

3: if jimg_loadsj¹/₄0 then

4: if jfailed_resolutionsj¹/₄1 then

5: Report this loading as a resolution failure

6: end if

7: else

8: spec Target Spec(S)

9: dirs DirSearchOrder(S)

10: if spec is the filename specification then

11: resolved path img_loads [0].resolved path // retrieve the first load

12: if IsUnsafeResolution (spec, resolved_path, dirs) then

```
13: Report this loading as an unsafe resolution
14: end if
15: end if
16: chained loads Chained Loading(S)
17: for each load in chained loads do
18: OfflineProfileAnalysis (each load)
19: end for
20: end if
```

IV. Discussion and Conclusion

In this paper, we have described the analysis technique to detect unsafe dynamic component loadings in networked systems by using IP address. Our technique works in two phases. It first generates profiles to record a sequence of component loading behaviours at runtime using dynamic binary instrumentation. It then analyzes the profiles to detect two types of unsafe component loadings: resolution failures and unsafe resolutions. To assess our technique, we implemented tools to detect unsafe component loadings on Microsoft Windows and Linux. Our evaluation shows that unsafe component loadings are prevalent on both platforms and more severe on Windows platforms from a security perspective.

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